

BPF/netdev implementation upgrade?

Agenda

01 bpf_sk_storage_get performance

02 What should we do about the cgroup's
bpf_sock_ops?

bpf_sk_storage_get()

≡

```
1829 bpf_sk_storage_get+0x64  
      8 bpf_sk_storage_get+0x68  
5075 bpf_sk_storage_get+0x6c
```

Two instructions are hot from the perf profile

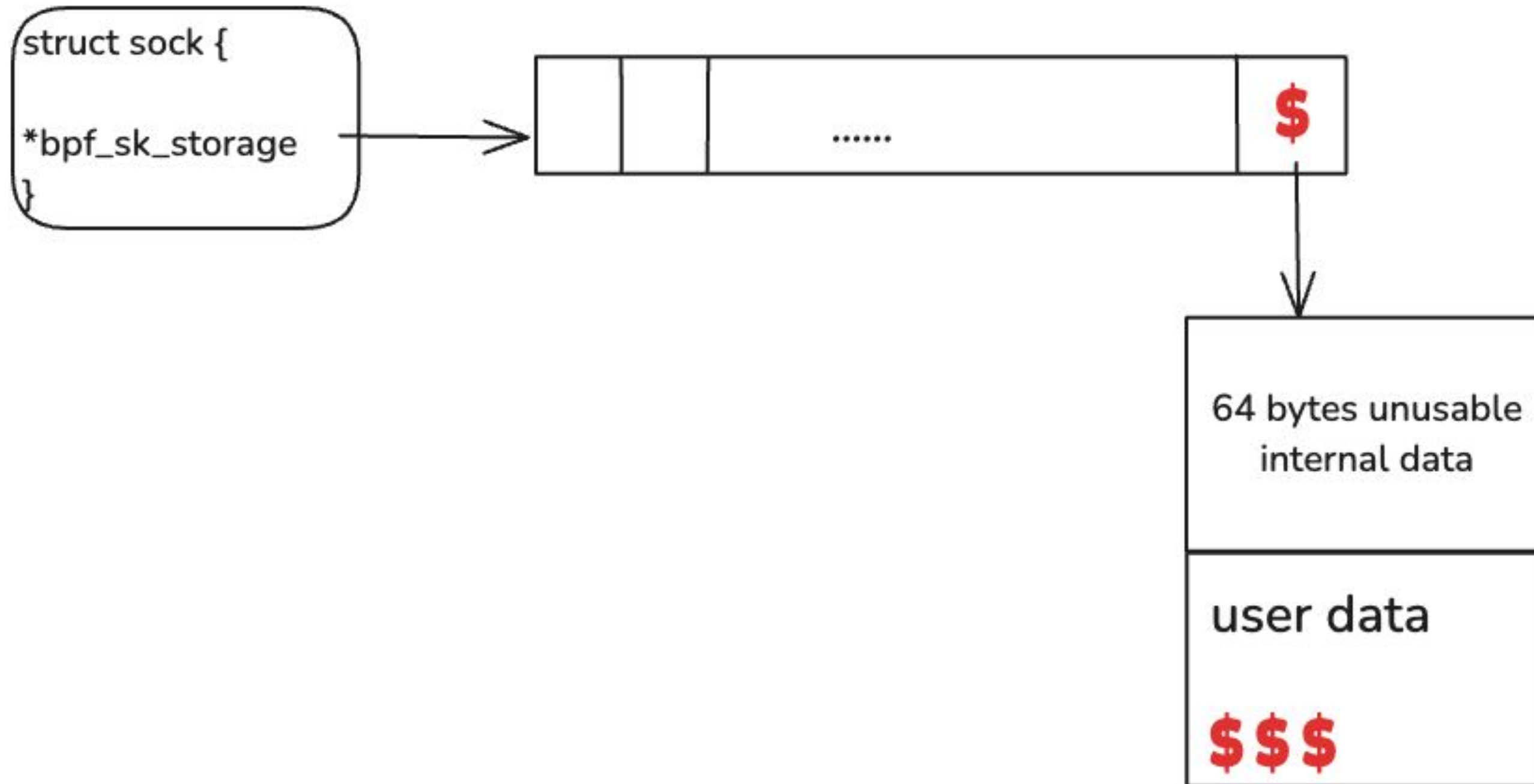
bpf_sk_storage_get()

```
static inline struct bpf_local_storage_data *
bpf_local_storage_lookup(struct bpf_local_storage *local_storage,
                        struct bpf_local_storage_map *smap,
                        bool cacheit_lockit)
{
    /* 1st cache-miss ($) */
    sdata = rcu_dereference_check(local_storage->cache[smap->cache_idx],
                                bpf_rcu_lock_held());

    /* 2nd cache-miss ($$$) */
    if (sdata && rcu_access_pointer(sdata->smap) == smap)
        return sdata;

    /* ... */
}
```

bpf_sk_storage_get()



Production usage observations

- The map_value's struct hardly changed
- In one of our maps, it had only one change since 2021
- The storage is needed for all sk(s) in a machine

What if...

```
struct sock {
```

```
    u8 sk_bpf_data[CONFIG_BPF_SIZE];
```

```
    struct fg_pkt_qos_info {  
        u32 marking_policy_id;  
        u32 egress_ent_policy_id;  
        /* ... */  
    };  
    struct flow_info {  
        char task[16];  
        u64 pid;  
        bool is_server;  
        /* ... */  
    };  
    .  
    .  
    .  
    .  
    .  
};
```

```
}
```

```
struct sock {
```

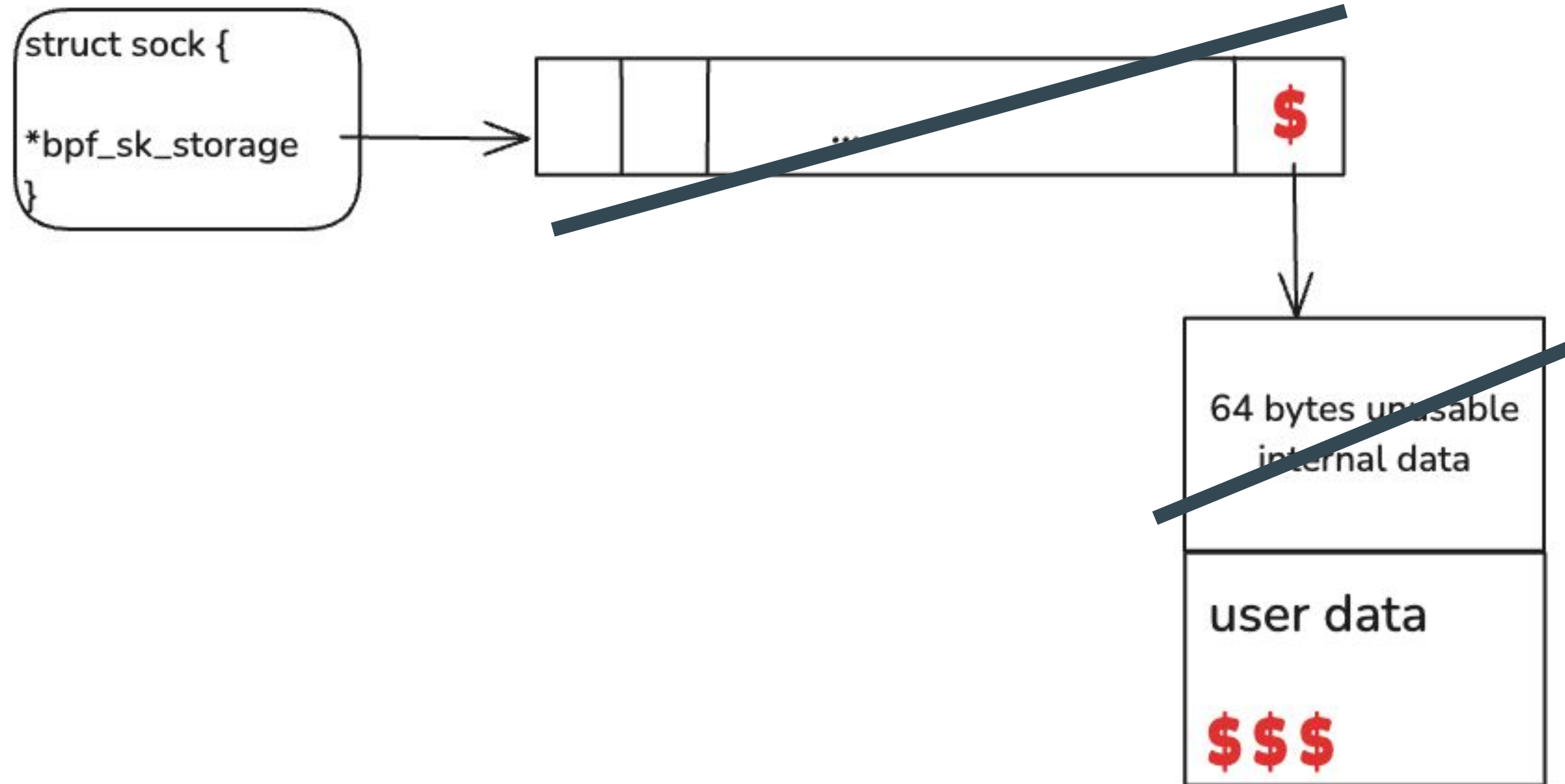
```
    u8 sk_bpf_data[CONFIG_BPF_SIZE];
```

```
    u32 marking_policy_id;  
    u32 egress_ent_policy_id;  
    char task[16]; /* ??? */  
    u64 pid;  
    bool is_server;
```

```
    .  
    .  
    .  
    .  
    .
```

```
}
```

What if...



Ideas...

- When CONFIG_BPF_SIZE is 0, goes back to existing slower path
- cgroup memory accounting is loss.

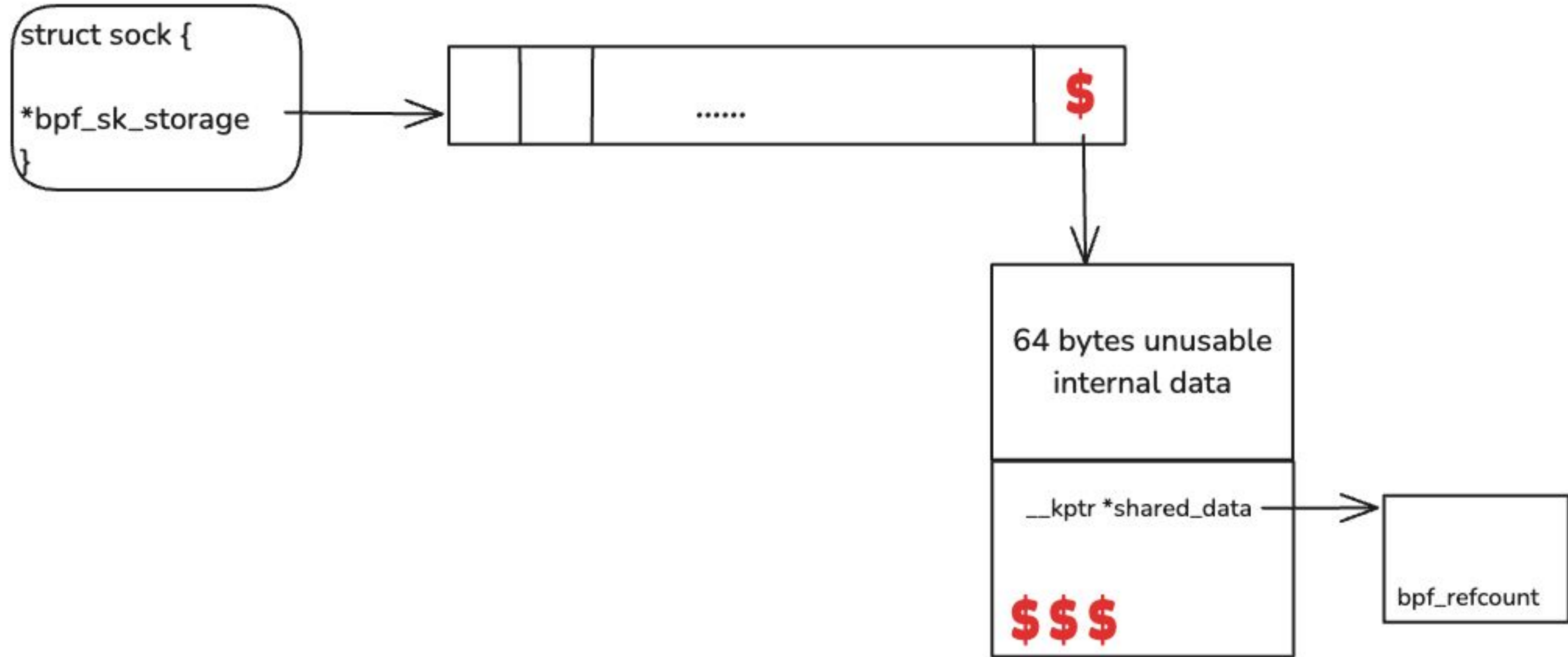
Ideas...

- How to reserve space in the sk_bpf_data?
 - First come first reserve.
 - Once reserved, the data will always be there until the reboot.
- The first 64 bytes internal data is gone.
 - The storage will stay with the sock(s) forever

Sharing sk storage data

- Half a million sockets having individual sk storages.
- The storages at most have 500 different values.
- Only 500 copies of the values are needed instead of half a million values.
- There is a `bpf_refcount` and `bpf_obj_new()`, so it should help?

Sharing sk storage data



What should we do with bpf_sock_ops?

```
enum {
    BPF_SOCKET_OPS_VOID,
    BPF_SOCKET_OPS_TIMEOUT_INIT,
    BPF_SOCKET_OPS_RWND_INIT,
    BPF_SOCKET_OPS_TCP_CONNECT_CB,
    BPF_SOCKET_OPS_ACTIVE_ESTABLISHED_CB,
    BPF_SOCKET_OPS_PASSIVE_ESTABLISHED_CB,
    BPF_SOCKET_OPS_NEEDS_ECN,
    BPF_SOCKET_OPS_BASE_RTT,
    BPF_SOCKET_OPS_RTO_CB,
    BPF_SOCKET_OPS_RETRANS_CB,
    BPF_SOCKET_OPS_STATE_CB,
    BPF_SOCKET_OPS_TCP_LISTEN_CB,
    BPF_SOCKET_OPS_RTT_CB,
    BPF_SOCKET_OPS_PARSE_HDR_OPT_CB,
    BPF_SOCKET_OPS_HDR_OPT_LEN_CB,
    BPF_SOCKET_OPS_WRITE_HDR_OPT_CB,
    BPF_SOCKET_OPS_TSTAMP_SCHED_CB,
    BPF_SOCKET_OPS_TSTAMP_SND_SW_CB,
    BPF_SOCKET_OPS_TSTAMP_SND_HW_CB,
    BPF_SOCKET_OPS_TSTAMP_ACK_CB,
    BPF_SOCKET_OPS_TSTAMP_SENDMSG_CB,
};
```

```
struct bpf_tcp_ops {
    u32 timeout_init(struct tcp_sock *sk);
    u32 rwnd_init(struct tcp_sock *sk);
    /* ... */
    u32 rtt_cb(struct tcp_sock *tp);
    u32 parse_hdr(struct tcp_sock *tp, struct sk_buff *skb);
    /* ... */
    U32 write_hdr_opt(struct tcp_sock *tp, struct sk_buff *skb);
};

struct bpf_sock_ops {
    void tx_tstamp(int tstamp_type, struct sk_buff *skb);
    void rx_tstamp(int tstamp_type, struct sk_buff *skb);
};
```

struct bpf_sock_ops_kern

```
struct bpf_sock_ops_kern {
    struct sock *sk;
    union {
        u32 args[4];
        u32 reply;
        u32 replylong[4];
    };
    struct sk_buff *syn_skb;
    struct sk_buff *skb;
    void *skb_data_end;
    u8  op;
    u8  is_fullsock;
    u8  is_locked_tcp_sock;
    u8  remaining_opt_len;
    u64  temp;          /* temp and everything after is not
                        * initialized to 0 before calling
                        * the BPF program. New fields that
                        * should be initialized to 0 should
                        * be inserted before temp.
                        * temp is scratch storage used by
                        * sock_ops_convert_ctx_access
                        * as temporary storage of a register.
                        */
};
```

Recent timestamp additions

	enum bpf_sock_ops	struct_ops
helpers/kfuncs filtering	runtime by skops->op. The op enum is a runtime value.	Verification time filtering by individual “ops” bpf prog.
kernel ptr access (e.g. sk)	is_valid_access() not useful because op enum is a runtime value. => BPF_INSN rewrite.	BTF. ops is known at the load time.
passing data to bpf prog	A kitchen sink in bpf_sock_ops_kern for all ops	Pass what is needed to the function’s arguments.
return values	Most ops ignore the CGRP_OK (1) and CGRP_EPERM (0). u32 replylong[4]. Not intuitive for bpf prog to use. union with the input args[4]. technically broken when >1 cgrp program running.	The function’s return value. Intuitive to bpf prog.

Discussions...

Q: Is the cgroup specific return value useful?

- The CGRP_OK (1) and CGRP_EPERM (0) are not very useful in sock_ops.
- Most of the ops ignore them. Equivalent to a struct_ops program returning void.
- Whatever done (e.g. bpf_setsockopt, writer_header...) is done. CGRP_OK/EPERM does not matter.

Q: Does it still need ordering?

- Should it be per-cgroup?
- or a global struct_ops ordering infra works for other sub-systems?

Q: Worth moving to struct_ops?

